



## Surviving in a world of change Towards evolvable and self-adaptive service-oriented systems

#### **ICSOC 2013**

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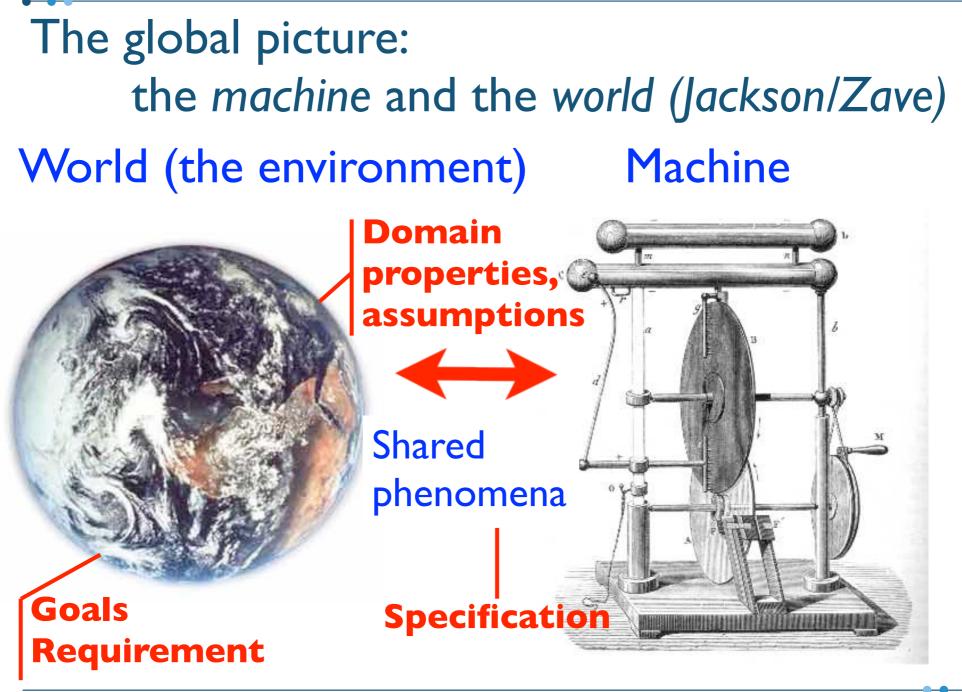
## Talk's roadmap

- Understanding change
- Understanding evolution and adaptation
- How can we detect change?
- How can we detect the need to evolve/adapt?
- How can we react to support evolution/adaptation?
- Lessons learned beyond SASs



## The root of the problems: endemic change





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## Domain properties and assumptions

- Both refer to problem world phenomena
- Properties hold regardless of any software-to-be
  - if a positive net force is applied in one direction then the body accelerates in that direction
  - (if plane has touched down then wheels turn)
- Assumptions may be violated
  - submission rate of user requests does not exceed XXX/sec
  - temperature is in the range -40 +40 Celsius
  - librarians register return of books when users bring borrowed books back



## **Domain assumptions**

May concern

- usage profiles
- users' responsiveness
- remote servers response time
- network latency
- sensors/actuators behaviors

"Domain assumptions bridge the gap between requirements and specifications" (M. Jackson & P. Zave)





## Dependability arguments

- Assume you have a formal representation for
  - -R = requirements
  - -S = specification
  - $-D = D_p + D_a$  domain properties and assumptions

if S and D are both satisfied and consistent, it is necessary to prove

– S, D |= R



## Change

- Requirements change
- Environment changes
- Change is often a manifestation of uncertainty
- Change asks for evolution (of the machine)



## Changes may cause evolution

- Changes are exogenous phenomena that may concern
  - R
  - D (actually,  $D_a$ )



- Changes likely break the dependability argument
- Evolution (of the machine) is a consequence of change
  - we need to change S (and hence the implementation) to continue to satisfy the dependability argument



### **Evolution and adaptation**

Adaptation is a special case of evolution due to changes in domain assumptions,  $D_a$ 

- an increasingly relevant phenomenon, often due to uncertainty
  - cyber-physical systems
    - interaction with the physical environment
  - user-intensive systems
    - changes in usage profile
  - cloud/service infrastructure
    - platform volatility

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Our focus here

## On-line vs off-line evolution (type vs instance) vs self-adaptive systems

- Traditionally, response to change is performed off-line by engineers (aka software maintenance)
- More and more often systems are required to be continuously running
- This asks for on-line evolution, i.e. applying changes to the machine as the system is running and providing service
- The special case of self-adaptive systems
  - (instance-level) self-managed on-line adaptation



## Self-adaptive system (SaS)

- D decomposed into  $D_f$  and  $D_c$ 
  - $D_f$  is the fixed/stable part
  - $D_c$  is the changeable part
- A SaS should
  - detect changes to D<sub>c</sub>
  - modify itself (the *machine* --- S, and the implementation) to keep satisfying the dependability argument, if necessary





## Paradigm shift

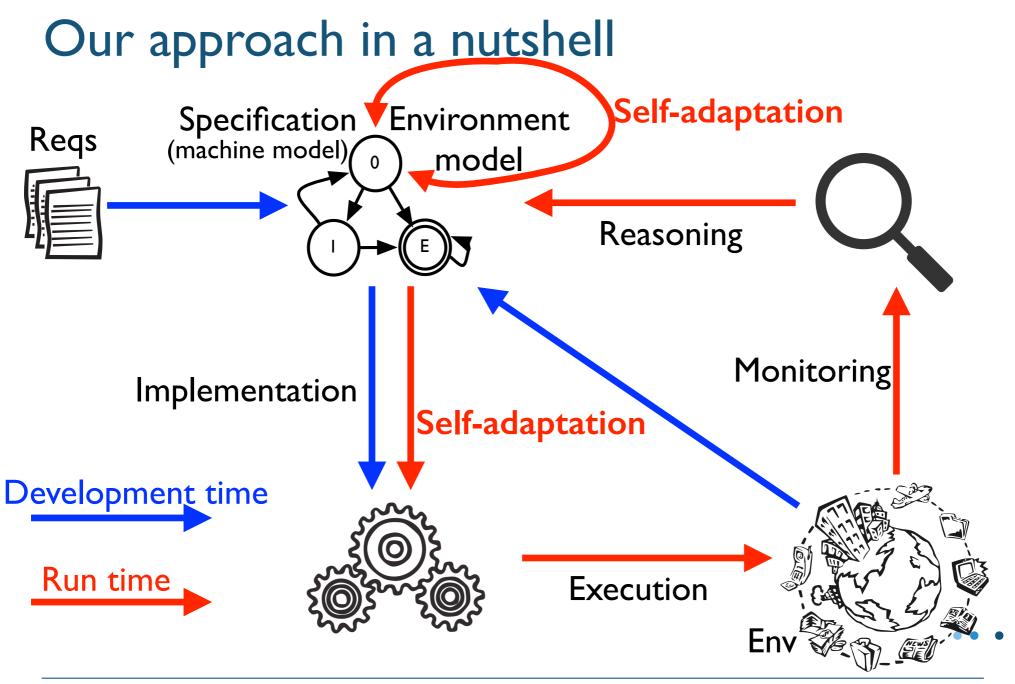
- SaSs ask for a paradigm shift, which involves both development time (DT) and run time (RT)
- The boundary between DT and RT fades
- Reasoning and reacting capabilities must enrich the RT environment
  - detect change
  - reason about themselves and the possible consequences of change
  - react to change



## Models+verification@runtime

- To detect change, we need to **monitor** the environment
- The changes must be retrofitted to models of the machine+environment that support reasoning about the dependability argument (a **learning** step)
- The updated models must be **verified** to check for violations to the dependability argument
- In case of a violation, a self-adaptation must be triggered









- I. Epifani, C. Ghezzi, R. Mirandola, G. Tamburrelli, "Model Evolution by Run-Time Parameter Adaptation", ICSE 2009
- C. Ghezzi, G. Tamburrelli, "Reasoning on Non Functional Requirements for Integrated Services", RE 2009
- I. Epifani, C. Ghezzi, G. Tamburrelli, "Change-Point Detection for Black-Box Services", FSE 2010
- A. Filieri, C. Ghezzi, G. Tamburrelli, "A formal approach to adaptive software: continuous assurance of non-functional requirements", Formal Aspects of Computing, 24, 2, March 2012.

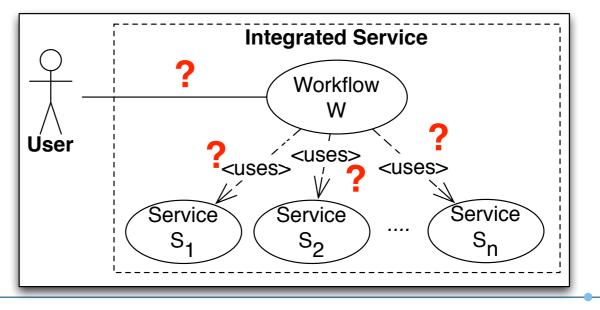


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## Zooming in

- Focus on non-functional requirements

   reliability, performance, energy consumption, cost, …
- Quantitatively stated in probabilistic terms
- $D_c$  decomposed into  $D_u$ ,  $D_s$ 
  - D<sub>u</sub> = usage profile
  - D<sub>s</sub> = S<sub>1</sub>  $\wedge$  ....  $\wedge$  S<sub>n</sub> S<sub>i</sub> assumption on i-th service



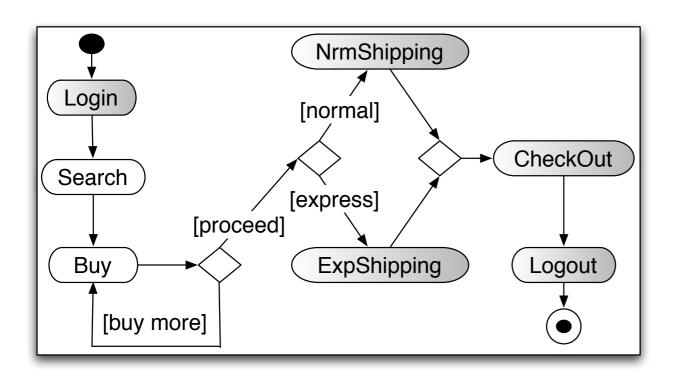


## Models

- Different models provide different viewpoints from which a system can be analyzed
- Focus on **non-functional** properties and quantitative ways to deal with uncertainty
- Use of Markov models
  - DTMCs for reliability
  - Reward DTMCs for energy/cost/performance..
- Use of probabilistic model checking for verification that a model satisfies a given property
  - Properties written in PCTL



## An example



Returning customers vs new customers

- 3 probabilistic requirements:
- RI: "Probability of success is > 0.8"
- R2:"Probability of a ExpShipping failure for a user recognized as ReturningCustomer < 0.035"

R3: "Probability of an authentication failure is less then < 0.06"

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## Assumptions

User profile domain knowledge

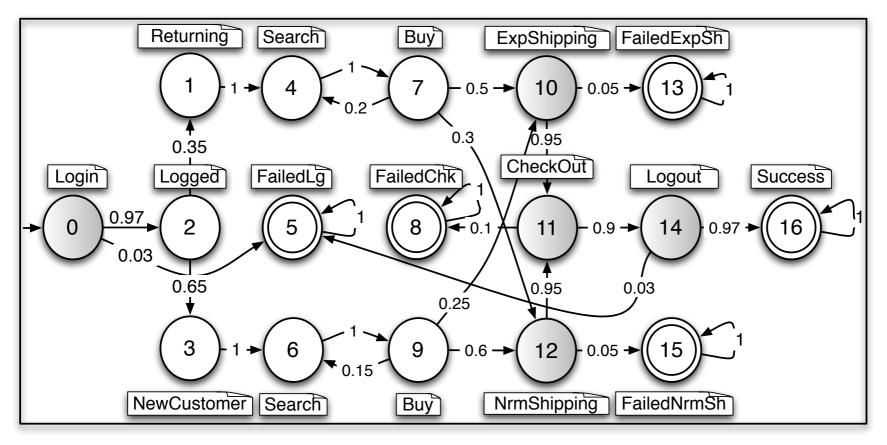
$D_{u,n}$	Description	Value
$D_{u,1}$	P(User is a RC)	0.35
$D_{u,2}$	P(RC chooses express shipping)	0.5
$D_{u,3}$	P(NCchooses express shipping)	0.25
$D_{u,4}$	P(RC searches again after a buy operation)	0.2
$D_{u,5}$	P(NC searches again after a buy operation)	0.15

#### External service assumptions (reliability)

$D_{s,n}$	Description	Value
$D_{s,1}$	P(Login)	0.03
$D_{s,2}$	P(Logout)	0.03
$D_{s,3}$	P(NrmShipping)	0.05
$D_{s,4}$	P(ExpShipping)	0.05
$D_{s,5}$	P(CheckOut)	0.1



## DTMC model



RI: "Probability of success is > 0.8" 0.84

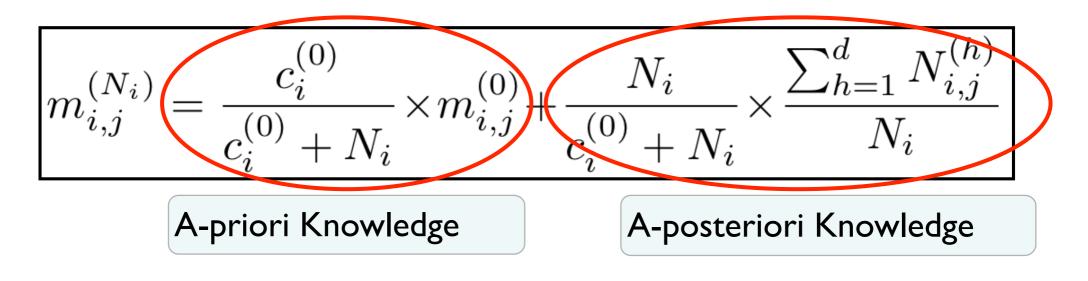
R2: "Probability of a ExpShipping failure for a user recognized as ReturningCustomer < 0.035" 0.031

R3: "Probability of an authentication failure is less then < 0.06" 0.056



## What happens at run time?

- Actual environment behavior is monitored
- Model updated by using a Bayesian approach to estimate DTMC matrix (posterior) given run time traces and prior transitions
- Boils down to the following updating rule





## Model update and failure prediction

- Model checking applied to after each update
- Model checking may predict requirements violations
- ... and trigger self-adaptations before violations manifest themselves



#### In our example





## The problem

- Verification subject to (application-dependent) hard real-time requirements
- Running conventional model checking tools after any change impractical in most realistic cases
- But changes are often local, they do not disrupt the entire specification
- Can they be handled in an **incremental** fashion?
- This requires revisiting model checking algorithms!



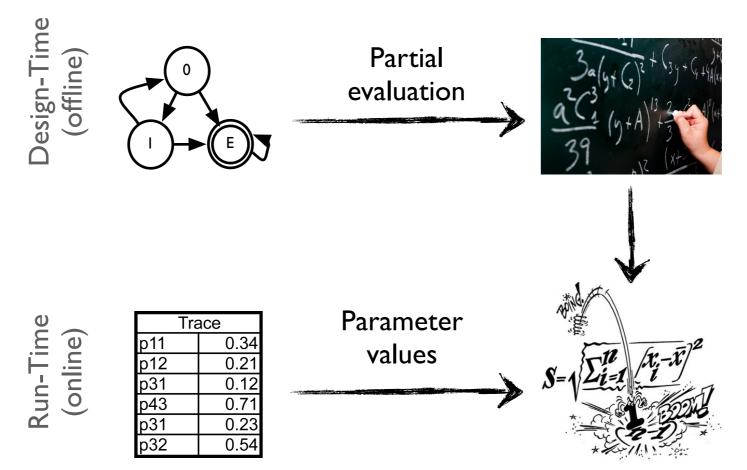
## Incrementality by parameterization

- Requires anticipation of changing parameters
- The model is partly numeric and partly symbolic
- Evaluation of the verification condition requires partial evaluation (mixed numerical/symbolic processing)
- Result is a formula (colynomial or reachability on DTMCs)

• Evaluation at run time substitutes actual values to symbolic parameters and is very efficient



## Working mom paradigm



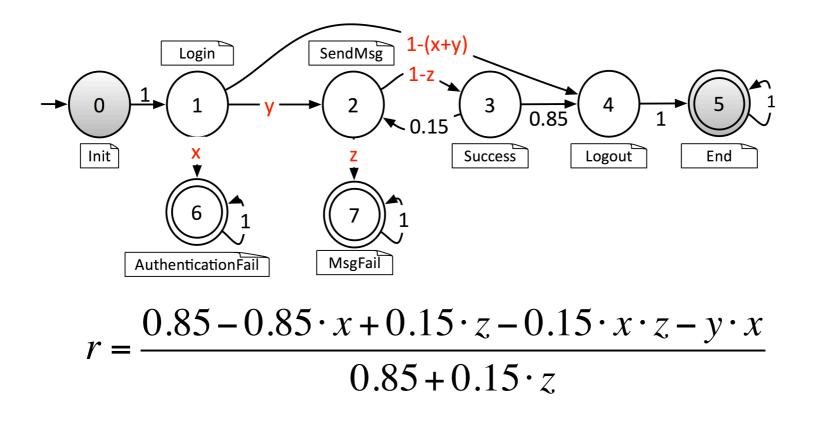
Analyzable properties: reliability, costs (e.g., energy consumption)

[ICSE 2011] A. Filieri, C. Ghezzi, G. Tamburrelli " Run-time efficient probabilistic model checking" [FormSERA 2012] A. Filieri, C. Ghezzi, "Further steps towards efficient runtime verification: Handling probabilistic cost models"

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## An example

$$r = \Pr(\diamondsuit{s} = 5) > r$$



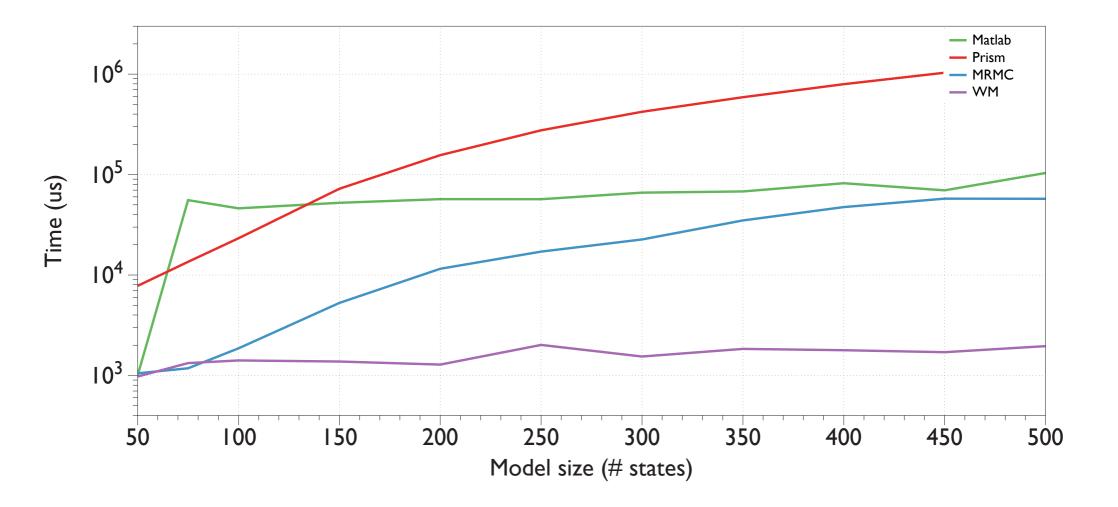


## The WM approach

- Assumes that the Markov model contains absorbing states, and that they are reachable
- Works by symbolic/numeric matrix manipulation
- All of (R) PCTL covered
- Expensive design-time partial evaluation, fast runtime verification
  - symbolic matrix multiplications, but very sparse and normally only few variables



#### Run-time verification



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### Further advantage of WM

- Because reachability properties can be expressed via polynomial functions, it is also possible to compute their (partial) derivative and perform sensitivity analysis
  - Which parameters affect most the global quality in the current operation point?
- Similar approach can deal also with rewards
  - Energy consumption, Average Execution time, Outsourcing cost, CPU time, Bandwidth



## The rest of the story

- After you detect the need for an adaptation, how do you react?
- You need to perform a dynamic update
- This means disconnecting components and ensuring a correct + safe update
- ... but this is subject for another talk



## What did we learn?

## How/where do we proceed from here?



## Run-time management

- The run-time environment for self-adaptive software should not *just run* applications
  - it should support introspection and reaction
    - on the application's requirements
    - ▶ its behaviour
    - the environment's behavior
- Models and continuous verification are essential for introspection and reaction
- But because models change, verification must be efficient

 $\checkmark$  constrained by real-time requirements

• This is agility taken to extremes



## Beyond self-adaptation

- Lessons learned are far reaching
  - Agile (explorative, incremental) development may become verification-driven by supporting incremental modelling and verification
  - Agility and formal methods may be reconciled rather than being antagonistic
- Vision
  - Towards verification-driven development as complementary to today's test-driven development



## Key feature: incrementality

#### **Incremental verification**

Given a system (model) S, and a set of properties P met by S Change = new pair S', P' where S'= S +  $\Delta$ S and P'= P +  $\Delta$ P

Let  $\prod$  be the proof of S against P The proof  $\prod$ ' of P' against S' can be done by just performing a proof increment  $\Delta \prod$  such that  $\prod' = \prod + \Delta \prod$ 

Expectations:

 $\Delta \Pi$  easy and efficient to perform  $\Delta \Pi$  helps designers reason about change

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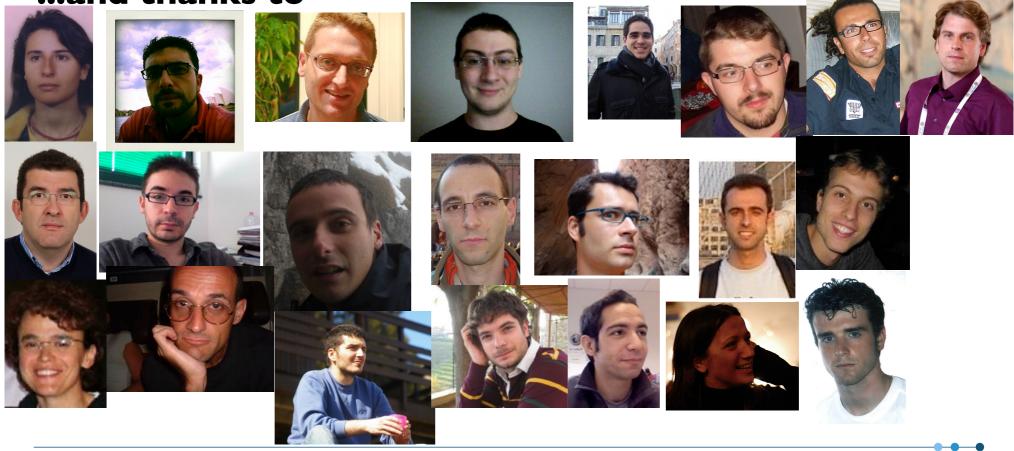
## A long way to go, but possible

- Revisit development models and verification procedures to make them incremental
- Make model-driven development practical
- Package above in IDEs



## Acknowledgements

- The work discussed here has been mostly developed thanks to a funding from the European Research Council (Advanced Grant IDEAS-ERC, Project 227977---SMScom)
- ...and thanks to



# Thanks for your attention

## **Questions?**